Tackling Neural Network Expressivity via (virtual Newton) Polytopes

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 - Can 3-layer NNs compute the maximum of 5 numbers?
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(both problems have broader context ...)

A Single ReLU Neuron



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Rectified linear unit (ReLU): $relu(x) = max\{0, x\}$



A Single ReLU Neuron



ReLU Feedforward Neural Networks

Acyclic (layered) digraph of ReLU neurons



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Acyclic (layered) digraph of ReLU neurons



Computes function

$$T_k \circ \operatorname{relu} \circ T_{k-1} \circ \cdots \circ T_2 \circ \operatorname{relu} \circ T_1$$

with linear transformations T_i .

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Acyclic (layered) digraph of ReLU neurons



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Example: depth 3 (2 hidden layers).

Theorem (Arora, Basu, Mianjy, Mukherjee (ICLR 2018)) $f: \mathbb{R}^n \to \mathbb{R}$ can be represented by a ReLU NN if and only if f is continuous and piecewise linear (CPWL). Support Functions and Newton Polytopes

Definition

The support function f_P of a polytope $P \subseteq \mathbb{R}^n$ maps a cost vector $x \in \mathbb{R}^n$ to the objective value of the linear program $\max_{v \in P} x^T v$.



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Definition

The *Newton polytope* of a convex CPWL function $f(x) = \max\{v_1^T x, v_2^T x, \dots, v_p^T x\}$ is $P(f) = \operatorname{conv}\{v_1, v_2, \dots, v_p\}$.

Bijection between Convex CPWL Functions and Polytopes

Have a bijection between these two sets:

- Convex CPWL functions $\mathbb{R}^n \to \mathbb{R}$,
- ▶ Polytopes in \mathbb{R}^n .

The bijection maps ...

- Functions to their Newton polytope $f \mapsto P_f$,
- Polytopes to their support function $P \mapsto f_P$.

Adding Two (Convex) CPWL Functions

$$\max\{0, x_1, x_2\} + \max\{0, x_2\} = \max\{0, x_2, x_1, x_1 + x_2, x_2, 2x_2\}$$
$$= \max\{0, x_1, 2x_2, x_1 + x_2\}$$

Adding Two (Convex) CPWL Functions

$$\max\{0, \mathbf{x_1}, \mathbf{x_2}\} + \max\{0, \mathbf{x_2}\} = \max\{0, \mathbf{x_2}, \mathbf{x_1}, \mathbf{x_1} + \mathbf{x_2}, \mathbf{x_2}, \mathbf{2x_2}\} \\ = \max\{0, \mathbf{x_1}, \mathbf{2x_2}, \mathbf{x_1} + \mathbf{x_2}\}$$



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▶ Minkowski sum: $P \otimes Q := \{p + q \mid p \in P, q \in Q\}$,

The Bijection is a Semigroup Isomorphism

$$P_{f+g} = P_f \otimes P_g,$$

$$f_{P \otimes Q} = f_P + f_Q.$$

$$\begin{array}{l} (\text{Convex CPWL functions } \mathbb{R}^n \to \mathbb{R}, \ +) \\ & \cong \\ (\text{Polytopes in } \mathbb{R}^n, \ \otimes) \end{array}$$

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- Can we do something about this?

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- Can we do something about this?

YES!

→ Virtual Polytopes!

► Semigroup (N, +)

Semigroup
$$(\mathbb{N}, +) \rightsquigarrow$$
 group $\mathbb{Z} = \{n - m \mid n, m \in \mathbb{N}\}$

- ▶ Semigroup $(\mathbb{N}, +) \rightsquigarrow$ group $\mathbb{Z} = \{n m \mid n, m \in \mathbb{N}\}$
- Semigroup $(\mathbb{Z} \setminus \{0\}, \cdot)$

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In the same way:

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Two Examples

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Cancellation Rule:

$$\left(\bullet \bullet \oslash \square \right) = \left(\bullet \oslash \right) = 1^{\otimes -1}$$

$$(\square \circ \square) \circ (\square \circ \checkmark)$$

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$$= (\nabla \otimes \varDelta) \otimes (\Box \otimes \checkmark) \\ (\nabla \otimes \Box) \otimes (\varDelta \otimes \checkmark)$$

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We even have a Group Isomorphism!

 $\begin{array}{l} (\mathsf{CPWL functions } \mathbb{R}^n \to \mathbb{R}, \ +) \\ \cong \\ (\mathsf{Virtual Polytopes in } \mathbb{R}^n, \ \otimes) \end{array}$

In particular:

Every CPWL function is a difference of two convex CPWL functions.

$$rac{P}{Q} \mapsto f_P - f_Q$$
 $f - g \mapsto rac{P_f}{P_g}$

Virtual Polytopes are Almost Polytopes

They have a face lattice,

faces are virtual polytopes.

They have a volume,

but volume can be negative.

Find out which virtual polytopes can occur as Newton polytopes of CPWL functions computed by neural networks of a certain size (or structure).

Addition

- Scalar multiplication
- ReLU activations

- Addition \leftrightarrow Minkowski addition \checkmark
- Scalar multiplication
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- Addition \leftrightarrow Minkowski addition \checkmark
- Scalar multiplication ?
- ReLU activations ?

We even have a Vector Space Isomorphism!

$$f = \max_{i=1}^{p} \{v_i^T x\} \quad \Rightarrow \quad P_{\lambda f} = \operatorname{conv}_{i=1}^{p} \{\lambda v_i\} = \lambda P_f$$

$\begin{array}{l} (\mathsf{CPWL functions, +, scalar multiplication}) \\ \cong \\ (\mathsf{Virtual Polytopes, } \otimes, \, \mathsf{scaling}) \end{array}$

(for $\lambda < 0$ also consistent with Minkowski difference)

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- ► Taking Maxima ?

For two convex CPWL functions f and g:

$$P_{\max\{f,g\}} = \operatorname{conv}\{P_f \cup P_g\} \eqqcolon P_f \oplus P_g.$$

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$$\Rightarrow \max\{f,g\} = \max\{f_+ - f_-, g_+ - g_-\} \\ = \max\{f_+ + g_-, g_+ + f_-\} - (f_- + g_-).$$

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Translating to virtual polytopes:

Does this remind you of something? $\rightsquigarrow \frac{a}{b} + \frac{c}{d} = \frac{ad+bc}{bd}$

We also have a Semiring Isomorphism!

 $\begin{array}{l} (\mathsf{CPWL functions, max, +}) \\ \cong \\ (\mathsf{Virtual Polytopes, } \oplus, \otimes) \end{array}$

(here you see the "tropical" world!)

- Addition \leftrightarrow Minkowski addition \checkmark
- ▶ Scalar multiplication \leftrightarrow Scaling of (virtual) polytopes \checkmark
- ► Taking Maxima ?

- Addition \leftrightarrow Minkowski addition \checkmark
- ▶ Scalar multiplication \leftrightarrow Scaling of (virtual) polytopes \checkmark
- Taking Maxima \leftrightarrow convex hull of the union \checkmark

Now to the two open problems:

- Can 3-layer NNs compute the maximum of 5 numbers?
- Can poly-size NNs solve the MST problem?

Computing the Maximum of Two Numbers

$$\max\{x, y\} = \max\{x - y, 0\} + y$$



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Computing the Maximum of Four Numbers



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lnductively: Maximum of *n* numbers with depth $\lceil \log_2(n) \rceil + 1$.

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Question: Is this best possible?

What's known?



• $\max\{0, x_1, x_2\}$ cannot be computed with 2 layers.

That's all!

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That's all!

- ▶ No function known that provably needs more than 3 layers.
- Smallest open case: Can max{0, x₁, x₂, x₃, x₄} be computed with 3 layers?

Virtual Newton Polytopes of Neural Networks



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 $\mathcal{P}_1 = \{ Y \oplus Z \mid Y, Z \text{ (virtual) zonotopes } \}$ (convex hull of the union of two zonotopes)



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(convex hull of the union of two zonotopes) $\mathcal{P}_2 = \{ \text{ finite Minkowski sums of (virtual) polytopes in } \mathcal{P}_1 \}$



 $\mathcal{P}_{1} = \{ Y \oplus Z \mid Y, Z \text{ (virtual) zonotopes } \}$ (convex hull of the union of two zonotopes) $\mathcal{P}_{2} = \{ \text{ finite Minkowski sums of (virtual) polytopes in } \mathcal{P}_{1} \}$ $f(x) = \max\{0, x_{1}, x_{2}, x_{3}, x_{4}\} \text{ can be computed by 3-layer NN}$ $\Leftrightarrow 4\text{-simplex } \Delta^{4} = P_{f} \in \mathcal{P}_{2}.$

The Two Open Problems

- Can 3-layer NNs compute the maximum of 5 numbers?
- Can poly-size NNs solve the MST problem?

The Minimum Spanning Tree Problem

$$G = \square$$

$$\mathcal{X} = \left\{ \square \right\}$$

Let $\mathcal{X} \subseteq \{0,1\}^E$ be the set of characteristic vectors of spanning trees in a fixed graph G = (V, E).

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MST problem: Compute the CPWL function

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MST problem: Compute the CPWL function

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Question: Is there a poly-size NN computing this function?

Is there a poly-size NN to solve MST?

$$f(c) = \min_{x \in \mathcal{X}} x^{\mathsf{T}} c = -\max_{x \in \mathcal{X}} (-x)^{\mathsf{T}} c$$

$$\Rightarrow \qquad P_f = \operatorname{conv} \{-x \mid x \in \mathcal{X}\}^{\otimes -1} = (-\mathsf{MST-Polytope})^{\otimes -1}$$

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Proposition

There exists a poly-size NN computing f

\Leftrightarrow

The MST-Polytope can be (virtually) generated from points via polynomially many operations of the form

Minkowski addition,

Convex hull of the union,

Scaling (with possibly inverting).

Thank you!



Questions? Ideas?