Computing the Maximum Function with ReLU Neural Networks

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A Single (Hidden) ReLU Neuron



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Rectified linear unit (ReLU): relu(x) = max{0, x}



ReLU Feedforward Neural Networks

Acyclic (layered) digraph of ReLU neurons



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ReLU Feedforward Neural Networks

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- Computes function T_k o relu o T_{k-1} o · · · o T₂ o relu o T₁ with affine transformations T_i.
- Example: depth 3 (2 hidden layers), width 3.

Example: Computing the Maximum of Two Numbers

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lnductively: Maximum of *n* numbers with depth $\lceil \log_2(n) \rceil + 1$.

Theorem (Wang, Sun (2005))

For any continuous and piecewise linear (CPWL) function $f : \mathbb{R}^n \to \mathbb{R}$, there are $\alpha_i \in \mathbb{R}$, $a_{ij} \in \mathbb{R}^n$, and $b_{ij} \in \mathbb{R}$ such that

$$f(x) = \sum_{i} \alpha_i \max\{a_{ij}^T x + b_{ij} \mid j = 1, \dots, n+1\}.$$

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 \Rightarrow Everything depends on the maximum function!

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- Smallest open case: Can max{0, x₁, x₂, x₃, x₄} be computed with 3 layers?
- ▶ No function known that provably needs more than 3 layers.

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(for notational purposes: $x_0 := 0$.)

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Construct a linear program (LP) that is feasible if and only if such a neural network exists.

Rough Idea



• f positively homogeneous if $f(\alpha x) = \alpha f(x)$ for all $\alpha > 0$.

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 \Rightarrow From now on, only consider NNs without biases.

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The $\binom{5}{2} = 10$ hyperplanes $x_i = x_j$, $0 \le i < j \le 4$, subdivide \mathbb{R}^4 into 5! = 120 cells in each of which the output of each neuron is affine.





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Claim: The output of a neuron in the first hidden layer is a linear combination of the following 14 functions:

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\Rightarrow 14 neurons suffice!

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Recall: cell = set of inputs with fixed ordering of 0, x_1, \ldots, x_4 .

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$$\Rightarrow$$
 2¹²¹ neurons suffice!

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 \rightsquigarrow For this talk suppose $2^{121}\approx 200\,000.$

Variables of the Linear Program



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Example: Cell $x_1 \ge x_2 \ge 0 \ge x_3 \ge x_4$: (1,0,0,0), (1,1,0,0), (0,0,-1,-1), (0,0,0,-1).

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- \Rightarrow Have these constraints:
 - 1. $4 \cdot 120$ equality constraints.
 - 2. $4 \cdot 120 \cdot 2^{121}$ inequality constraints.

... for less than a minute ...

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... and outputs ...

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Infeasible!

Three Obvious Next Steps ...

Proving the assumption.

► Finding a "real" proof.

Generalize to more layers.

Thank you!



Questions? Ideas?